Goroutines demystified.

Mathieu Nativel





A goroutine is a lightweight thread managed by the Go runtime.



A goroutine is a lightweight thread managed by the Go runtime.

They are called lightweight threads because they require less processing time :

- Smaller default stack size
- · Lighter context switching : setup and teardown don't require call to the kernel.



Ok so basically they are just user threads : an implementation of threads and scheduling running on top of the OS.



Ok so basically they are just user threads : an implementation of threads and scheduling running on top of the OS.





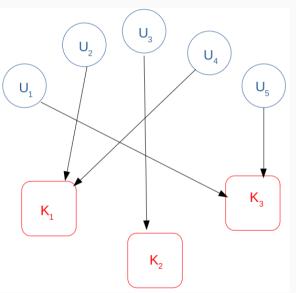
Some nice ideas to make profit of user threads :

- Allocate kernel threads when creating the first user threads.
- Park for reuse the kernel threads after the user thread ends.
- Schedule user threads to run on respecting kernel threads.



User-threads implementation

Multiplexing low-cost user-threads on high-cost kernel threads :





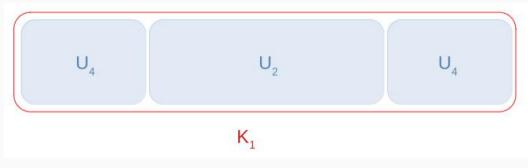


Figure 2: Multiple user thread can run on the same associated kernel thread.



The scheduler manages a runqueue of runnable goroutines.

When it wants to schedule a goroutine it pops it out of the runqueue and schedules it on a available kernel thread (instantiating one if needed and possible).



The scheduler manages a runqueue of runnable goroutines and per-core runqueues of runnable goroutines.

Work-stealing scheduling: When it wants to schedule a goroutine on a kernel thread it first try to pop it from the local runqueue, if failed it tries to steal it from other local runqueues and finally it tries the global runqueue.



Go implements theses concepts through 3 important structures in the runtime code :

The M struct (Kernel Thread)	The P struct (Linked List)
- Represents a kernel thread - Contains two important pointers: one to the currently running G and	- Represents a scheduling context - Contains a list of runnable Gs.
	Represents a kernel thread Contains two important pointers:



From the go runtime source code : https://golang.org/src/runtime/runtime2.go

```
type g struct {
stack
            stack // offset known to runtime/cgo
               unsafe.Pointer // passed parameter on wakeup
param
atomicstatus
               uint32
goid
               int64
schedlink
               guintptr
waitsince
                         // approx time when the g become blocked
               int64
waitreason
               waitReason // if status==Gwaiting
tracelastp
               puintptr // last P emitted an event for this goroutine
    };
```



type m struct {

curg	<pre>*g // current running goroutine</pre>
caughtsig	<pre>guintptr // goroutine running during fatal signal</pre>
р	<pre>puintptr // attached p for executing go code (nil if not executing go code)</pre>
nextp	puintptr
oldp	<pre>puintptr // the p that was attached before executing a syscall</pre>
id	int64
spinning	bool // m is out of work and is actively looking for work
};	

go

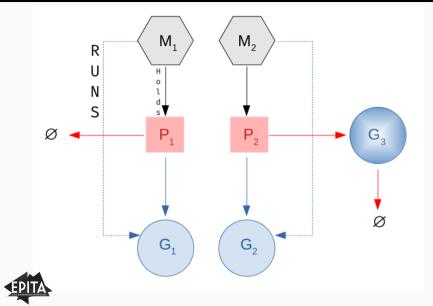


```
type p struct {
     id
                 int32
 status
            uint32 // one of pidle/prunning/...
             muintptr // back-link to associated m (nil if idle)
 m
 runghead uint32
 rungtail uint32
          [256]guintptr
 runa
 // runnext. if non-nil. is a runnable G that was ready'd by
 // the current G and should be run next instead of what's in
 // rung if there's time remaining in the running G's time
 // slice. It will inherit the time left in the current time
 // slice ...
 runnext guintptr
     . . .
```

};



Maybe with a drawing ?

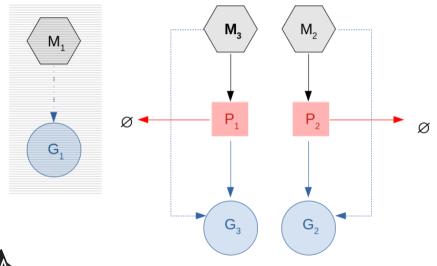


The idea is that when a M (a kernel thread) is blocked (on a syscall for exemple), the go scheduler can take its runqueue (*its P*) and give it to an other M.



What happens when blocking

Blocked on syscall





Questions ?



- author: LSE title: Presentation ...

